AUTOMATA 2010

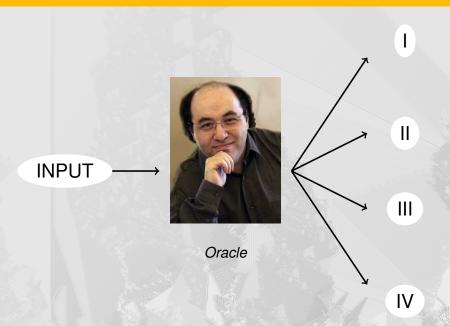
Guillaume Theyssier

LAMA lab. (CNRS, Université de Savoie, France)

June 14th, 2010

Chronology

- Banks (1970)
- Albert and Čulik (1987)
- Martin I (1993)
- Durand and Róka (1996)
- Durand-Lose (1997)
- Mazoyer and Rapaport (1998)
- Ollinger (2002)
- T (2005)



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 - 1 define a priori a list of properties or classes
 - 2 being similar ⇔ satisfying the same properties

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 - define a relation of similarity or equivalence between CA
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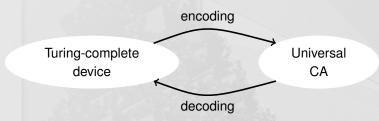
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Key points:

- no external tools used to define properties
- infinitely many classes

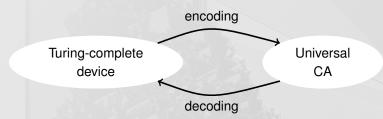
What is universality?

► Classical approach:



What is universality?

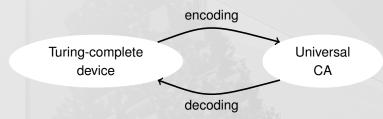
► Classical approach:



- encoding/decoding/halt problem
- only a positive definition

What is universality?

► Classical approach:



- encoding/decoding/halt problem
- only a positive definition
- ► This talk:
 - 1 define a notion of **simulation** between CA
 - 2 intrinsic universality ^{def} ability to simulate *all* CA

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 - define a **pre-order** on CA: <u>≺</u>

 (reflexive and transitive relation)

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 - 2 classification:
 - induced equivalence relation

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topology of the pre-order

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$$F \sim G \stackrel{def}{\Leftrightarrow} F \leq G \text{ and } G \leq F$$

- topology of the pre-order
- 3 universality:

F universal
$$\stackrel{def}{\Leftrightarrow} \forall G, G \leq F$$

Simulation pre-orders

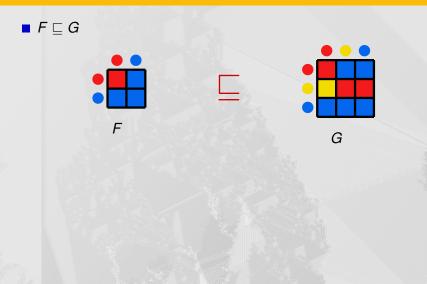
► Ingredients:

- 1 "local" comparison relations
 - subautomaton
 - factor
- rescaling operations

General idea

pre-order ≡ local comparison up to rescaling



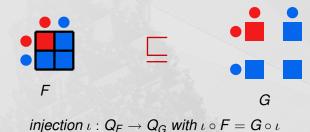


F ⊆ G

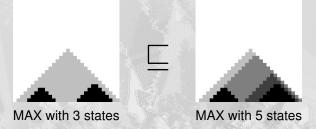


injection $\iota: Q_F \to Q_G$ with $\iota \circ F = G \circ \iota$

F ⊆ G



example:



Factor



Factor

■ *F* ⊴ *G*



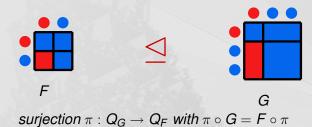
 \leq



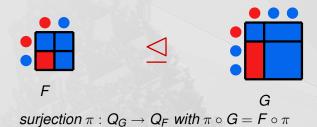
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Factor

F ⊴ G



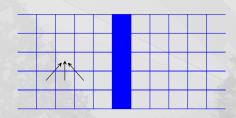
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example



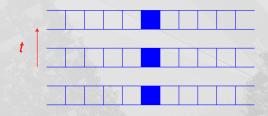
▶ 3 parameters: $F \mapsto F^{\langle m,t,z \rangle}$



Global map

$$F^{\langle 1,1,0 \rangle} = F$$

- ▶ 3 parameters: $F \mapsto F^{\langle m,t,z \rangle}$
- time

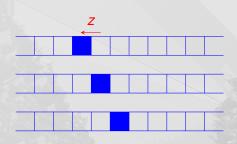


Global map

$$F^{\langle 1,t,0\rangle}=F^t$$

▶ 3 parameters: $F \mapsto F^{\langle m,t,z\rangle}$

- time
- shift

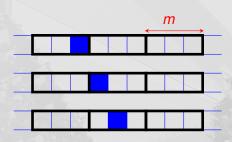


Global map

$$F^{\langle 1,t,z\rangle} = \sigma_z \circ F^t$$

▶ 3 parameters: $F \mapsto F^{\langle m,t,z\rangle}$

- time
- shift
- cell grouping



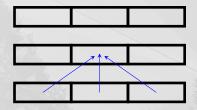
Global map

$$F^{\langle m,t,z\rangle} = \mathbf{0}_m^{-1} \circ \sigma_z \circ F^t \circ \mathbf{0}_m$$

 $\mathbf{o}_m^{-1}: Q^{\mathbb{Z}} \to (Q^m)^{\mathbb{Z}}$ canonical bijection

▶ 3 parameters: $F \mapsto F^{\langle m,t,z\rangle}$

- time
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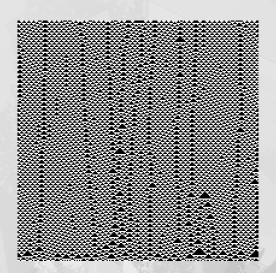


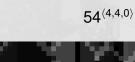
Fact

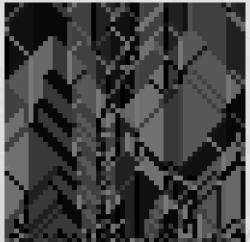
 $F^{\langle m,t,z\rangle}$ is a cellular automaton

- with a possibly different alphabet
- with a possibly different radius

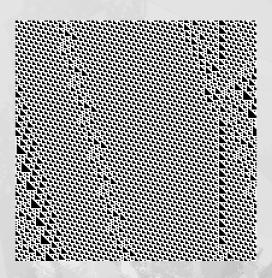
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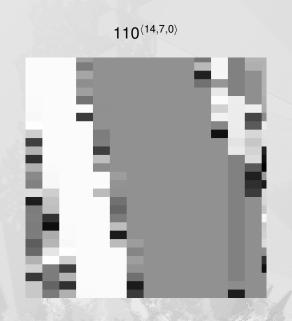






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3 pre-orders



3 pre-orders

■ injective simulation

$$F \preceq_{i} G \stackrel{def}{\Leftrightarrow} \exists \overrightarrow{p_{1}}, \overrightarrow{p_{2}} : F^{\langle \overrightarrow{p_{1}} \rangle} \sqsubseteq G^{\langle \overrightarrow{p_{2}} \rangle}$$

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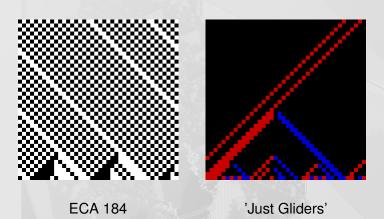
■ surjective simulation

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mixed simulation

$$F \leq_{\mathbf{m}} G \stackrel{def}{\Leftrightarrow} \exists \overrightarrow{p_1}, \overrightarrow{p_2} : F^{\langle \overrightarrow{p_1} \rangle} \trianglelefteq \sqsubseteq G^{\langle \overrightarrow{p_2} \rangle}$$

Examples

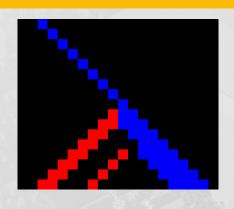




'Just Gliders' \leq 184 $^{\langle 2,2\rangle}$ \sim 184



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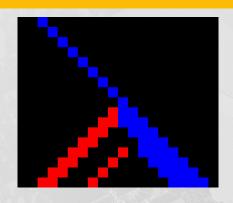


'Just Gliders'
$$\leq$$
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'Just Gliders' ≤_s 184



'Just Gliders' \sqsubseteq 184 $^{\langle 2,2\rangle}$ \sim 184







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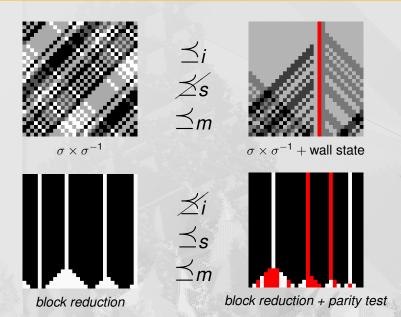
'Just Gliders' ≤, 184

Separation

Separation



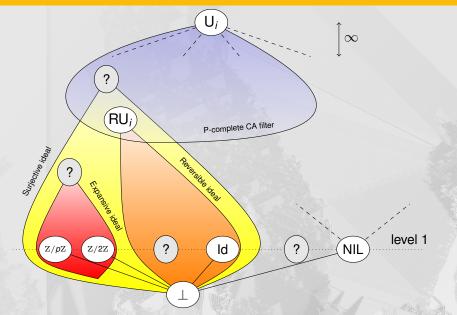
Separation



Pre-order *≤*_i

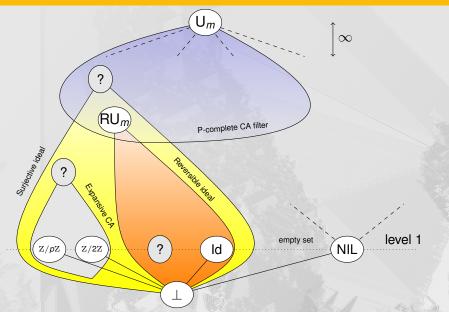


Pre-order \leq_i





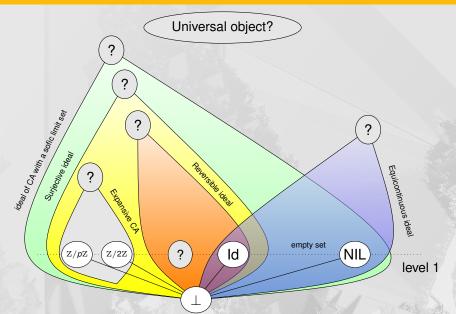
Pre-order \leq_m



Pre-order \leq_s



Pre-order \leq_s



intrinsic universality $\stackrel{def}{=} U_i \stackrel{def}{=} \{F : \forall G, G \leq_i F\}$



N. Ollinger, "Universalities in Cellular Automata", Handbook of natural computing

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Theorem (Rapaport) — no real-time universality

If $F \in U_i$ then there is some G with $G^{\langle \overrightarrow{p} \rangle} \not\sqsubseteq F^{\langle t,t,z \rangle}$, $(\forall \overrightarrow{p},t,z)$.

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Exercise

If $F \in U_i$ and has a spreading state then it can simulate any G without using the spreading state.

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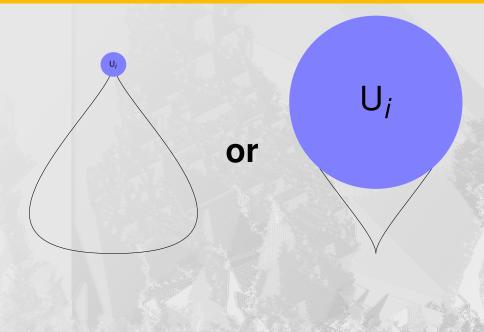
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Exercise

If F is not universal then there is a non-universal G with

$$F \leq_i G$$
 but $G \not \leq_i F$



The quest for small intrinsically universal CA



The quest for small intrinsically universal CA

- ▶ 2D
 - Banks 1970: 2 states + von Neuman neighb.
 - Conway 1970: Game of Life (2 states + Moore neighb.)

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- Ollinger 2002: 6 states + radius 1
- Richard 2008: 4 states + radius 1 (U_i or U_m ?)

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Open problems

- $U_i \stackrel{?}{=} U_m$
- is 110 intrinsically unievrsal?

Adding syntactical constraints



Adding syntactical constraints

- **▶** Definitions:
 - **F** is **captive** if $\forall \overrightarrow{x}$:

$$f(x_1,\ldots,x_k)\in\{x_1,\ldots,x_k\}$$

F is **multiset** if $\forall \overrightarrow{x}$ and \forall permutation π

$$f(\overrightarrow{x}) = f(\pi(\overrightarrow{x}))$$

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Theorem

There exists intrinsically universal CA in the following families:

- number conserving
- totalistic (⊆ multiset)
- captive

Density

- \blacksquare \mathcal{P} a **property** (a set of CA)
- \blacksquare \mathcal{F} a **family** (a set of CA)
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Theorem (Boyer,T.)

- $d_n(U_i/\text{captive CA}) = 1$
- $d_r(U_i/\text{multiset CA}) = 1$
- combination of both + other families (e.g. outer-totalistic)

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Open problem

- \blacksquare $d_n(U_i/CA)$?
- $d_r(U_i/CA)$?

syntactical restriction ⇒ restriction on global behaviours?

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Proposition

- \blacksquare $\forall F$ there is G with radius 1 and $F \sim G$
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▶ Encodings

- lacktriangledown $\phi: CA
 ightarrow \mathcal{F}$ encoding if $F \preceq_i \phi(F)$ for all F
- **fairness** wrt universality: $F \in U_i \Leftrightarrow \phi(F) \in U_i$

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Theorem

There are recursive fair encodings from CA into captive CA, and from CA into multiset CA.

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Open problem

What are equivalence classes without any captive and/or multiset CA?



- ► Reversible CA
 - there is no reversible $F \in U_i$

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- but "reversible-universality" (RU_i) exists!

Theorem (Durand-Lose)

There exists a reversible *F* with $G \leq_i F$ for all reversible *G*.

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▶ Surjective CA

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Open problem

Is there a "surjective-universal" CA in dimension 1?

- ▶ Pre-order \leq_s
 - **q** persistent state $\stackrel{def}{\Leftrightarrow} f(*,q,*) = q$
 - $X \stackrel{def}{=} CA$ with a persistent state

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Open problem

Is there a universal CA for \leq_s ?

Future directions

- 1 more general simulations?
 - sub-systems induced by stable sub-shifts (SFT? sofic?)
 - factor with context (sliding block codes)
- 2 dimension change
 - constructions à la Toffoli?
 - sub-actions à la Hochman?

3 asynchronous CA

Don't Know What to Do During the World Cup?

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▶ My list of small open questions

- **1** *F* non-surjective $\stackrel{?}{\Rightarrow}$ NIL $\leq_i F$
- **2** F reversible $\stackrel{?}{\Rightarrow}$ Id \leq F
- 3 find a reversible F with $F \nsim F^{-1}$
- 4 more than 3 classes of equicontinuous (up to σ) CA?
- **5** F expansive and $G \subseteq F \stackrel{?}{\Rightarrow} G$ expansive

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Thank you!